



William G. Kolakoski

Run-length encoding operator / fixed point

Herbert Freeman

Chain codes / uniformity / balance

Azriel Rosenfeld

Hierarchy of runs



The run-length encoding operator

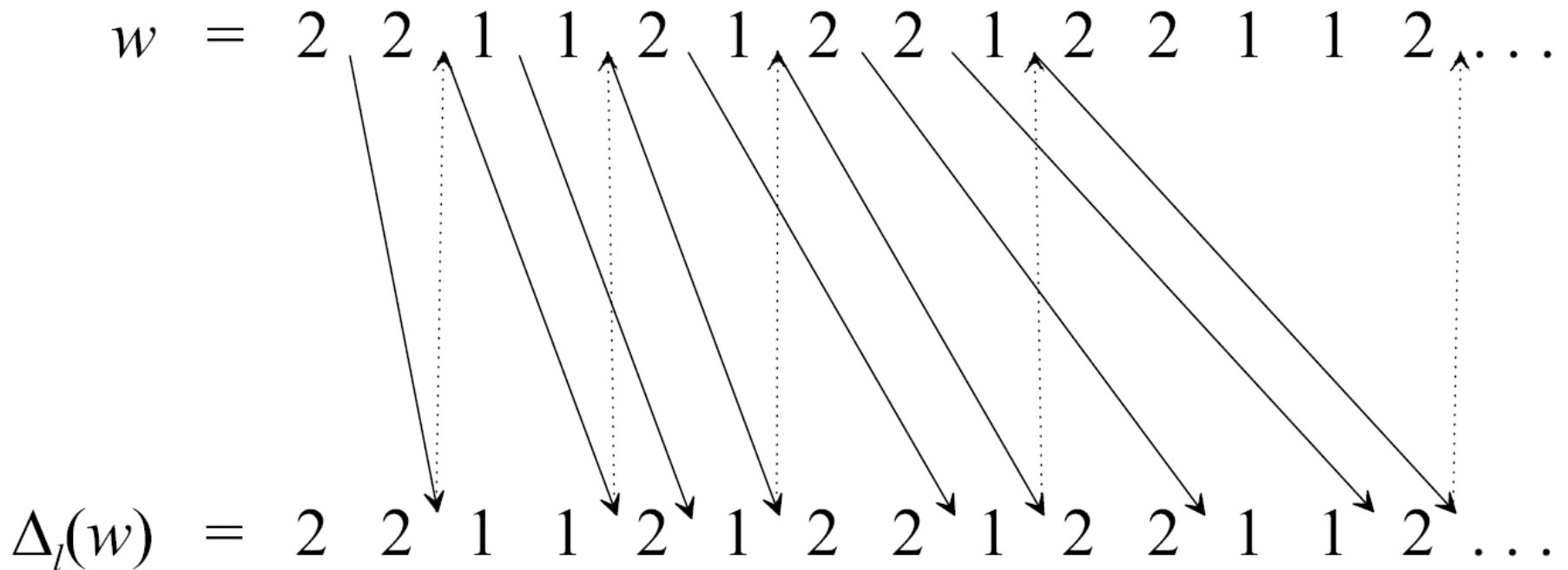
$$\Delta_l: \{1, 2\}^\omega \rightarrow \mathbf{N}^\omega$$

$$w = \begin{cases} 1^{k_1} 2^{k_2} 1^{k_3} 2^{k_4} \dots, & \text{if } w \in 1 \cdot \{1, 2\}^\omega \\ 2^{k_1} 1^{k_2} 2^{k_3} 1^{k_4} \dots, & \text{if } w \in 2 \cdot \{1, 2\}^\omega \end{cases}$$

$$\Delta_l(w) = k_1 k_2 k_3 \dots$$



The Kolakoski word





Digital lines...

$$D_{R'}(y = ax, x > 0) = \{(k, \lceil ak \rceil); k \in \mathbf{N}^+\}$$

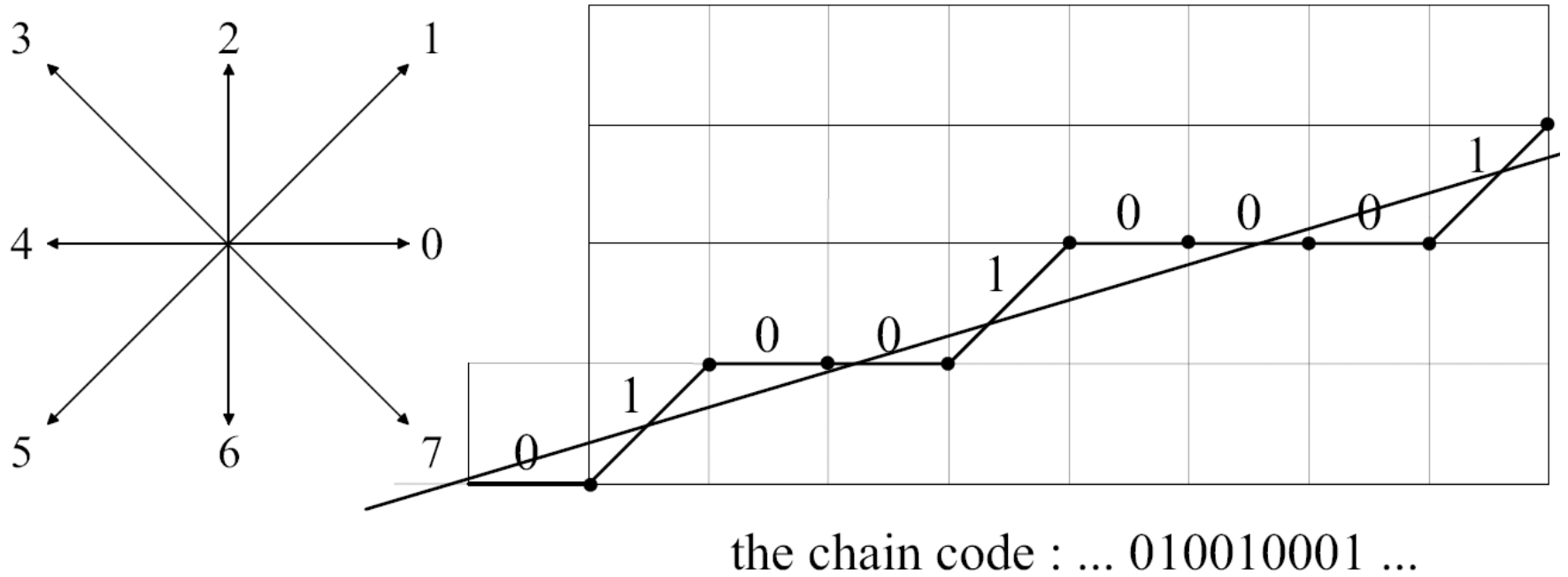
...and binary words

$$s'(a): \mathbf{N} \rightarrow \{0, 1\}$$

$$\forall n \in \mathbf{N} \quad s'_n(a) = \lceil a(n+1) \rceil - \lceil an \rceil$$



The Freeman code



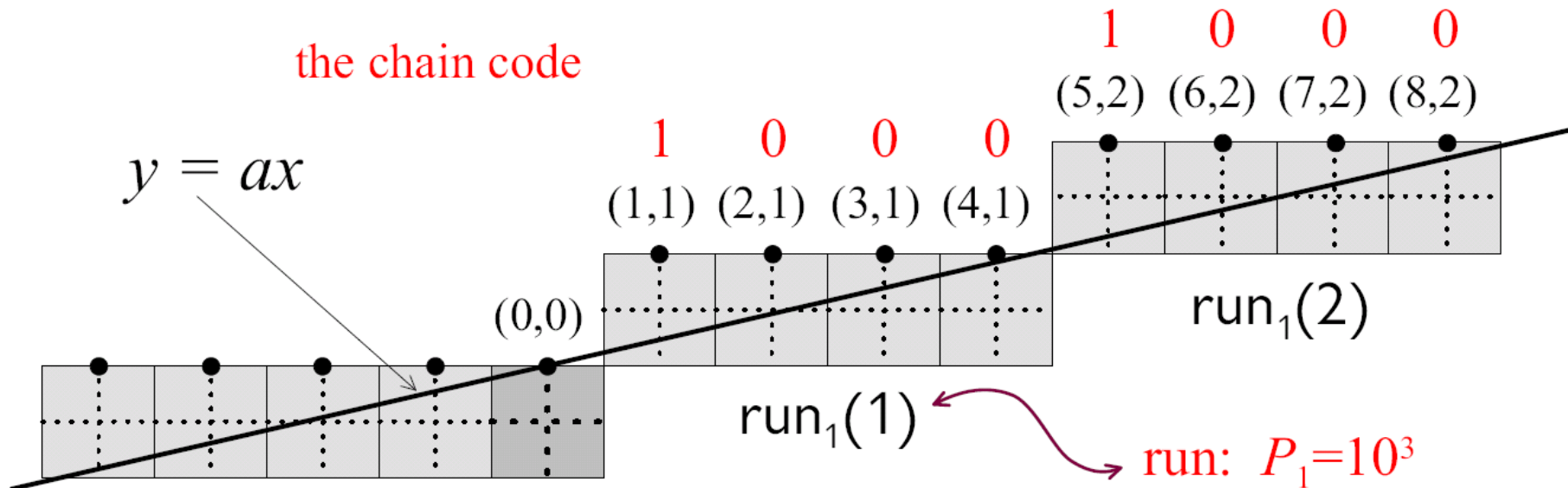


Freeman's notion of balance

(F3) successive occurrences of the element occurring singly are as uniformly spaced as possible.



Rosenfeld and his runs





Rosenfeld and his runs

Two run lengths on level 1: runs₁ $S_1=10^m$ and $L_1=10^{m+1}$

Two run lengths on level 2: runs₂ $S_2=S_1L_1^k$ and $L_2=S_1L_1^{k+1}$

or $S_2=S_1^kL_1$ and $L_2=S_1^{k+1}L_1$

Two run lengths on level n : runs _{n} $S_n=S_{n-1}L_{n-1}^l$ and $L_n=S_{n-1}L_{n-1}^{l+1}$

or $S_n=S_{n-1}^lL_{n-1}$ and $L_n=S_{n-1}^{l+1}L_{n-1}$

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Three questions. About:

the run length on level n

the main run on level $n-1$

the first run on level $n-1$



The index jump function

$$a = [0; a_1, a_2, a_3, a_4, a_5, a_6, a_7, \dots]$$

$$i_a : \mathbf{N}^+ \rightarrow \mathbf{N}^+$$

$$i_a(1) = 1, \quad i_a(2) = 2, \quad \text{for } n \geq 2:$$

$$i_a(n + 1) = i_a(n) + 1 + \delta_1(a_{i_a(n)})$$

$$\begin{array}{cccccccccccccccc}
 a = [0; & \overset{b_1}{\underline{1}}, & \overset{b_2}{a_2}, & \overset{b_3}{\underline{1}, 1}, & \overset{b_4}{a_5}, & \overset{b_5}{\underline{1}, 1}, & \overset{b_6}{a_8}, & \overset{b_7}{a_9}, & \overset{b_8}{\underline{1}, a_{11}}, & \overset{b_9}{a_{12}}, & \overset{b_{10}}{\underline{1}, 1}, & \overset{b_{11}}{\underline{1}, a_{16}}, & \overset{b_{12}}{a_{17}}, & \dots] \\
 & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \dots \\
 (i_a(k))_{k \in \mathbf{N}^+} = & (& 1, & 2, & 3, & 5, & 6, & 8, & 9, & 10, & 12, & 13, & 15, & 17, & \dots)
 \end{array}$$



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The slope of the line (upper mech. word) is $a = [0 ; a_1, a_2, a_3, \dots]$

How $i_a(k)$ and $a_{i_a(k)}$ describe the form of run_k :

$S_1 = 10^{a_1-1}$, $L_1 = 10^{a_1}$, and, for $k \geq 2$:

$$a_{i_a(k)} \geq 2 \quad \Rightarrow \quad S_k = S_{k-1}^{a_{i_a(k)}-1} L_{k-1} \quad (\text{or } L_{k-1} S_{k-1}^{a_{i_a(k)}-1}),$$

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Sturmian words with balanced construction

slopes: $a = [0; 1, 2, 1, 1, 3, 1, 1, a_8, a_9, \dots]$

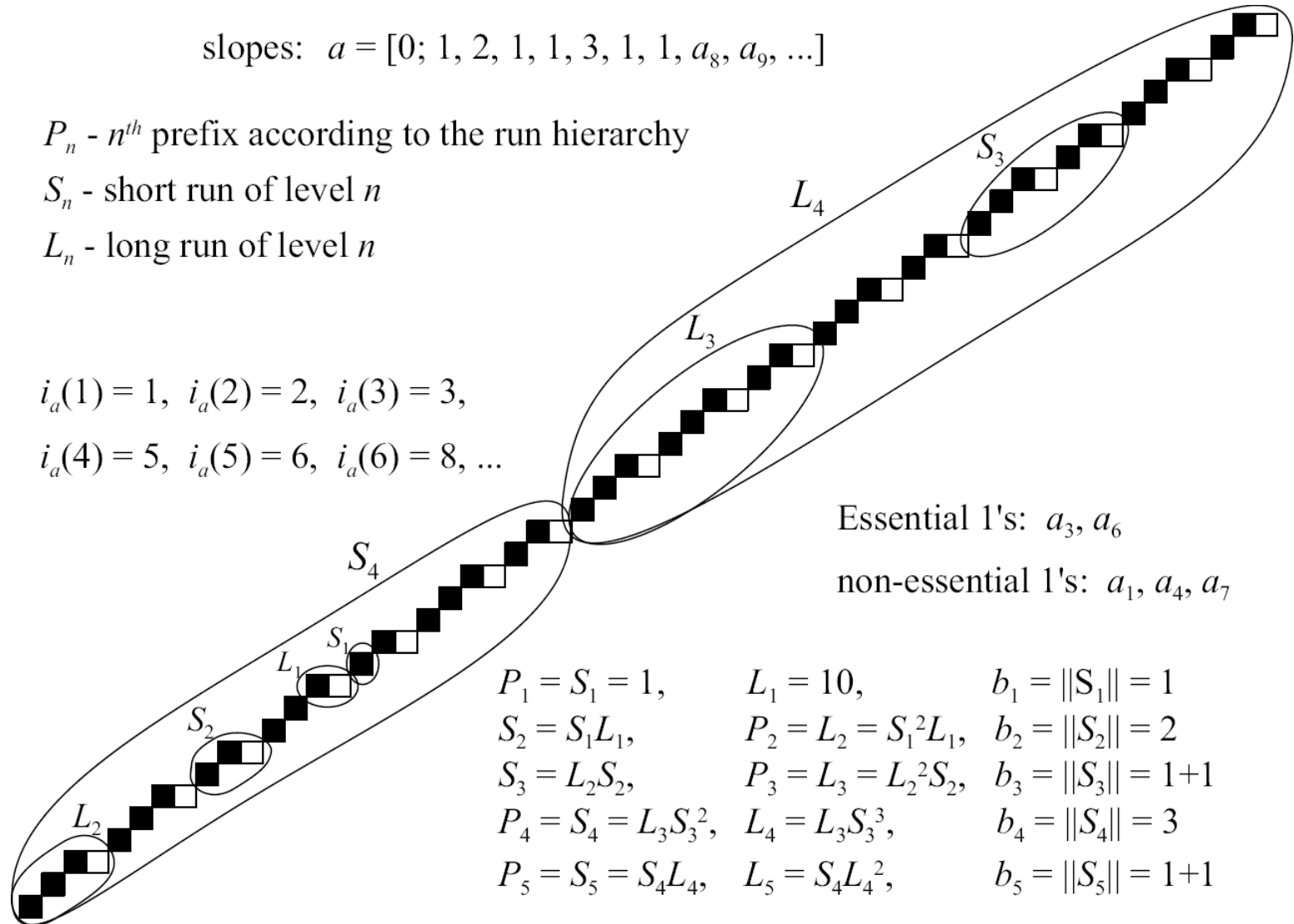
P_n - n^{th} prefix according to the run hierarchy

S_n - short run of level n

L_n - long run of level n

$i_a(1) = 1, i_a(2) = 2, i_a(3) = 3,$

$i_a(4) = 5, i_a(5) = 6, i_a(6) = 8, \dots$



Essential 1's: a_3, a_6

non-essential 1's: a_1, a_4, a_7

$P_1 = S_1 = 1,$	$L_1 = 10,$	$b_1 = \ S_1\ = 1$
$S_2 = S_1L_1,$	$P_2 = L_2 = S_1^2L_1,$	$b_2 = \ S_2\ = 2$
$S_3 = L_2S_2,$	$P_3 = L_3 = L_2^2S_2,$	$b_3 = \ S_3\ = 1+1$
$P_4 = S_4 = L_3S_3^2,$	$L_4 = L_3S_3^3,$	$b_4 = \ S_4\ = 3$
$P_5 = S_5 = S_4L_4,$	$L_5 = S_4L_4^2,$	$b_5 = \ S_5\ = 1+1$



Calculations for the latest example

k	$a_{i_a(k+1)} = 1?$	main_k	$i_a(k+1)$	b_k	prefix P_k of $s'(a)$
1	$a_{i_a(2)} = a_2 = 2 \geq 2$	S_1	even	1	$S_1 = 1$
2	$a_{i_a(3)} = a_3 = 1$	L_2	odd	2	$L_2 = S_1^2 L_1 = 1110$
3	$a_{i_a(4)} = a_5 = 3 \geq 2$	S_3	odd	2	$L_3 = L_2^2 S_2$
4	$a_{i_a(5)} = a_6 = 1$	L_4	even	3	$S_4 = L_3 S_3^2$
5	$a_{i_a(6)} = a_8 \geq 2$	S_5	even	2	$S_5 = S_4 L_4$
6	$a_{i_a(7)} = a_9 \geq 2$	S_6	odd	a_8	$L_6 = S_5^{a_8} L_5$



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5	$a_{i_a(6)} = a_8 \geq 2$	S_5	even	2	$S_5 = S_4 L_4$
6	$a_{i_a(7)} = a_9 \geq 2$	S_6	odd	a_8	$L_6 = S_5^{a_8} L_5$



Calculations for the latest example

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2	$a_{i_a(3)} = a_3 = 1$	L_2	odd	2	$L_2 = S_1^2 L_1 = 1110$
3	$a_{i_a(4)} = a_5 = 3 \geq 2$	S_3	odd	2	$L_3 = L_2^2 S_2$
4	$a_{i_a(5)} = a_6 = 1$	L_4	even	3	$S_4 = L_3 S_3^2$
5	$a_{i_a(6)} = a_8 \geq 2$	S_5	even	2	$S_5 = S_4 L_4$
6	$a_{i_a(7)} = a_9 \geq 2$	S_6	odd	a_8	$L_6 = S_5^{a_8} L_5$



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Two equivalence relations on the set of slopes

1. based on **run length** on all levels for $s'(a)$:

$$a \in [(b_1, b_2, b_3, \dots)] \sim_{\text{len}} \iff \forall k \in \mathbf{N}^+ \quad \|S_k\| = b_k$$

2. based on **run construction** on all levels for $s'(a)$:

$$a \sim_{\text{con}} a' \iff i_a \equiv i_{a'}$$



The constructional word $\gamma(a) \in \{0, 1\}^\omega$

Let $a = [0; a_1, a_2, \dots]$. For $n \in \mathbf{N}^+$:

$$\gamma_n(a) = i_a(n+2) - i_a(n+1) - 1$$

$$\gamma_n(a) = \delta_1(a_{i_a(n+1)})$$

$$\gamma_n(a) = \begin{cases} 0, & S_n \text{ is the most frequent} \\ & \text{run on level } n \text{ for } s'(a) \\ 1, & L_n \text{ is the most frequent} \\ & \text{run on level } n \text{ for } s'(a). \end{cases}$$



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Definition The *run-construction encoding operator* $\Delta_c : \mathcal{UM}_0 \longrightarrow \{0, 1\}^\omega$ is defined as $\Delta_c = (1\gamma) \circ (s')^{-1}$.

$$\begin{array}{ccc}]0, 1[\setminus \mathbf{Q} & \xrightarrow{s'} & \mathcal{UM}_0 \\ & \searrow 1\gamma & \downarrow \Delta_c \\ & & \{0, 1\}^\omega \supset \mathcal{UM}_0 \end{array}$$

where \mathcal{UM}_0 denotes the set of all upper mechanical words with irrational slope $0 < a < 1$ and with intercept 0.



Let $a \in]0, 1[\setminus \mathbf{Q}$. The word $s'(a) = 1c(a)$ has
balanced construction if

$$\exists \alpha \in \mathbf{R} \quad \gamma(a) = c(\alpha)$$

Sturmian-balanced construction if

$$\exists \alpha \in]0, 1[\setminus \mathbf{Q} \quad \gamma(a) = c(\alpha)$$

self-balanced construction

$$1\gamma(a) = \Delta_c(1c(a)) = 1c(a)$$



A fixed-point theorem:
exactly 1 fixed point in each equivalence class

Let $(b_n)_{n \in \mathbf{N}^+}$ be such that $b_1 \in \mathbf{N}^+$
and $b_n \in \mathbf{N}^+ \setminus \{1\}$ for all $n \geq 2$. Then

$$\exists^1_{a \in]0,1[\setminus \mathbf{Q}}$$

$$a \in [(b_n)_{n \in \mathbf{N}^+}]_{\sim_{\text{len}}} \wedge s'(a) = \Delta_c(s'(a)).$$



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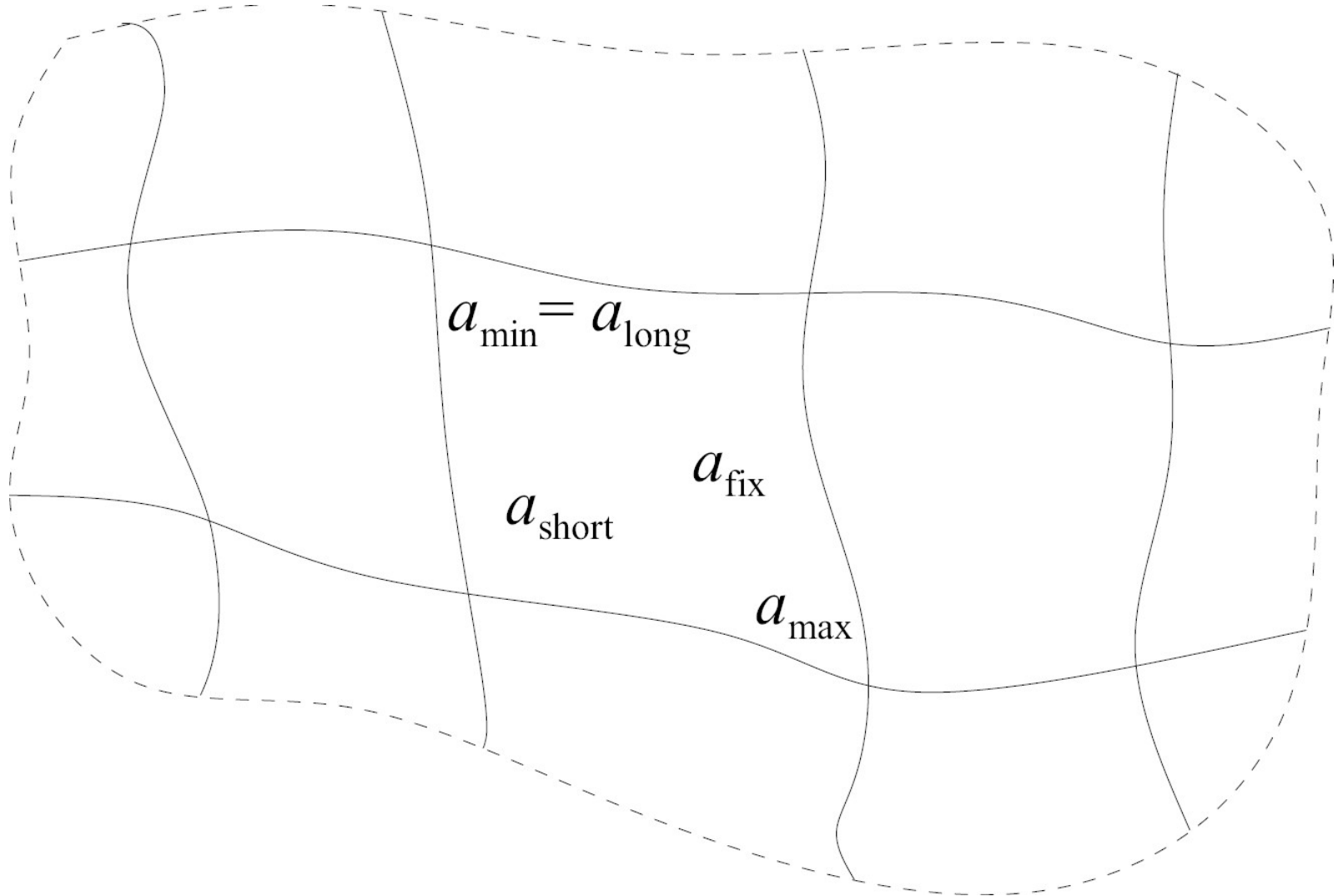
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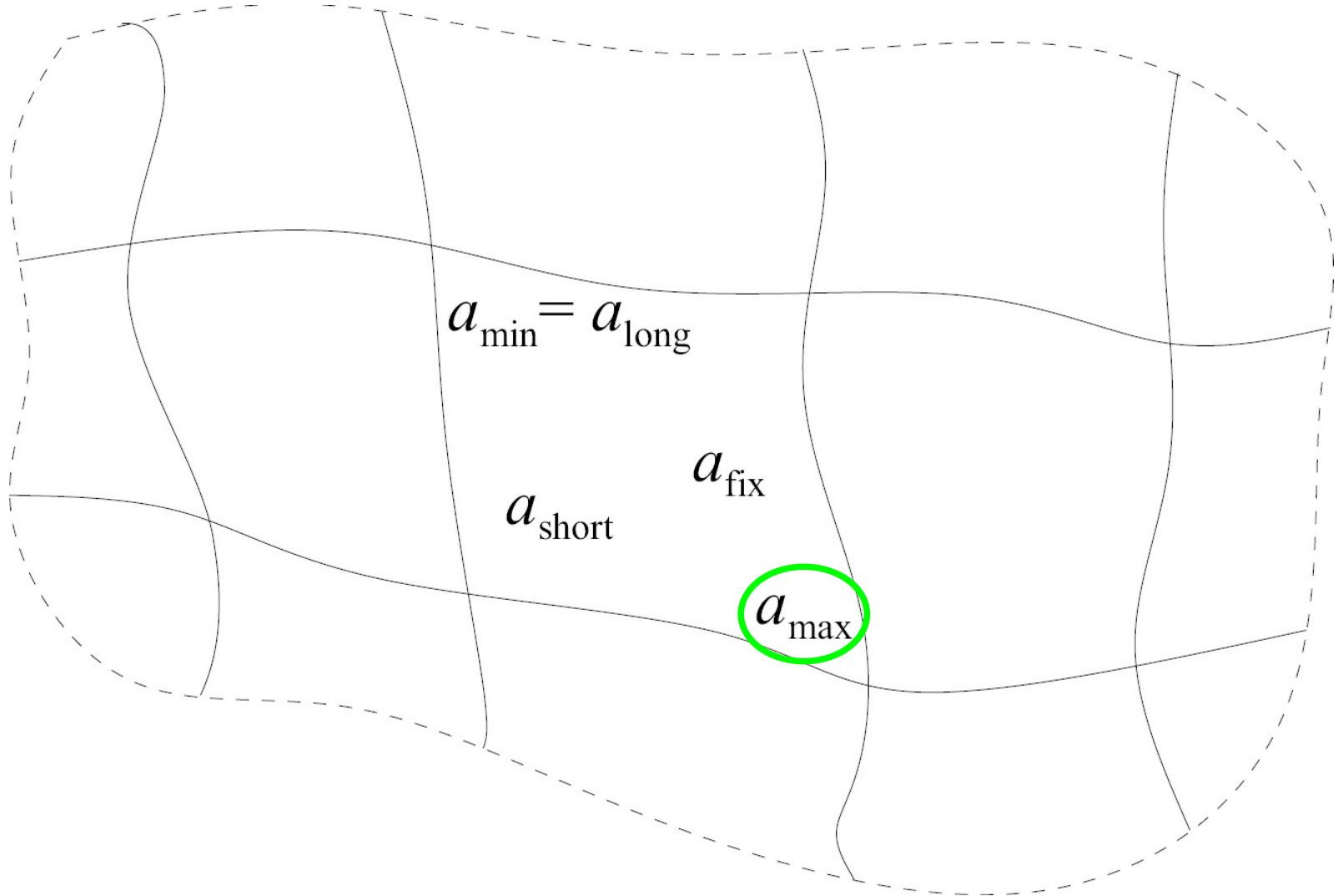


Sturmian words with balanced construction



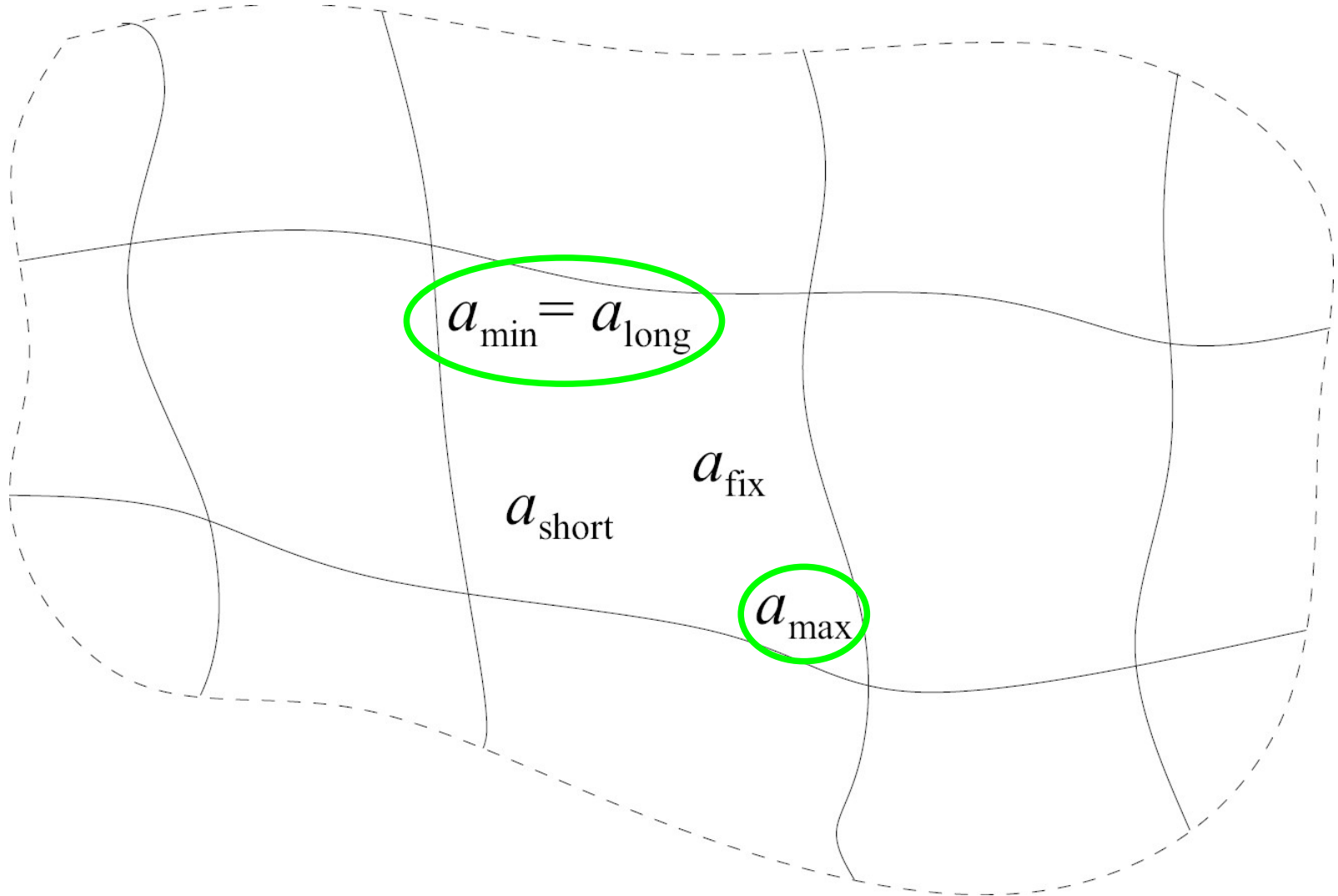


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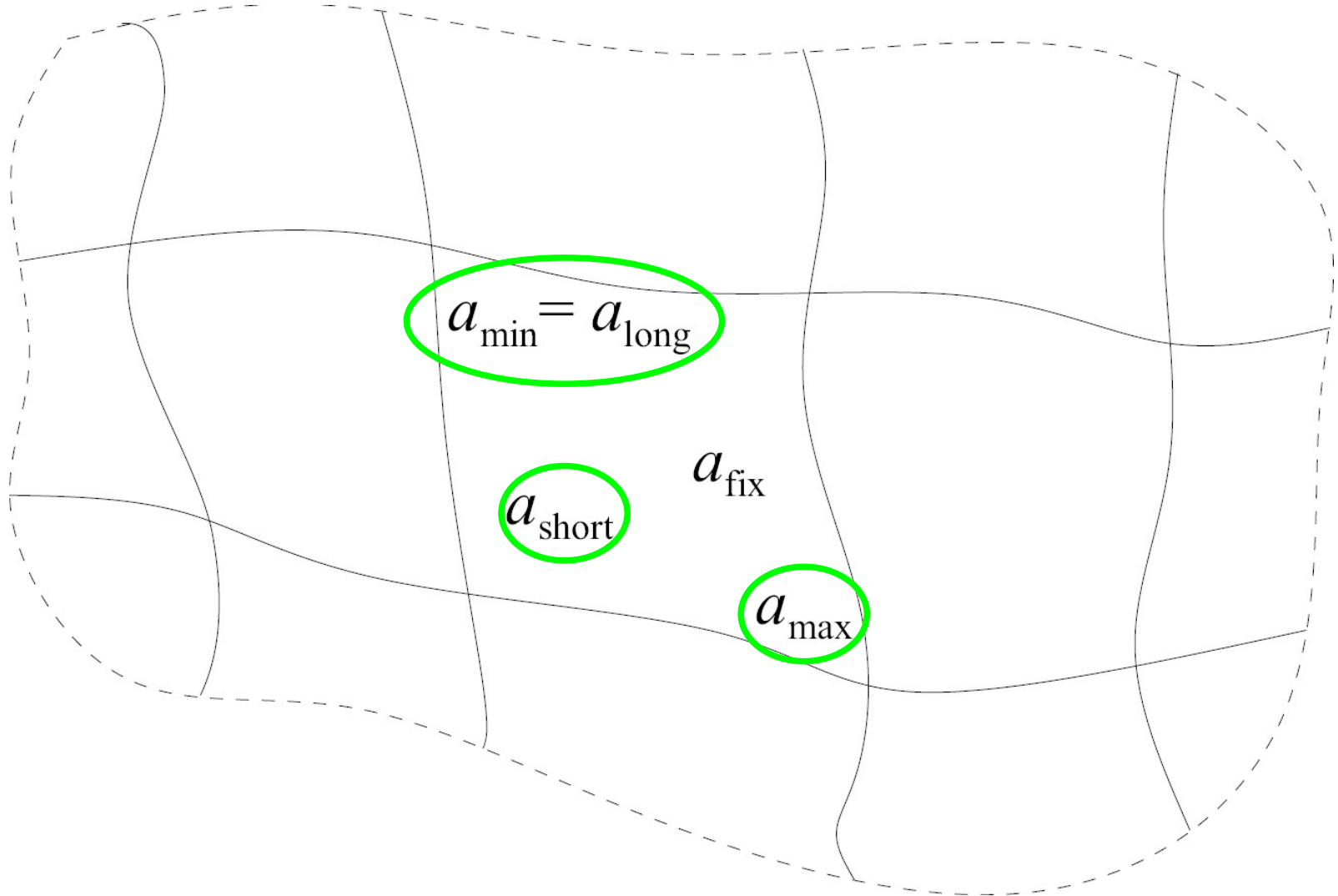


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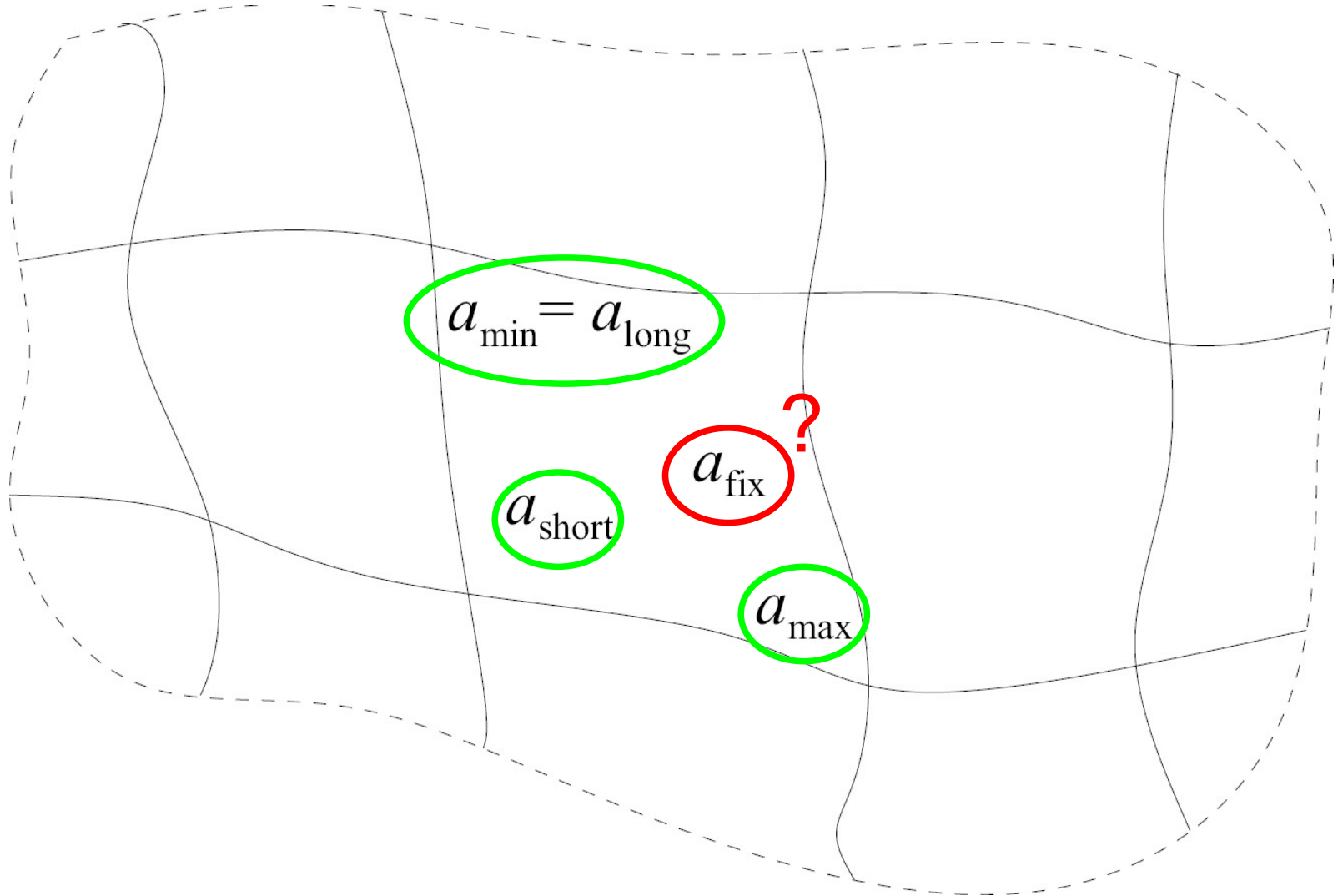


Sturmian words with balanced construction





Sturmian words with balanced construction





Theorem Let $\text{Fix}(\Delta_c) \subset \mathcal{UM}_0$ denote the set of all fixed points of Δ_c . Then:

1. $\text{Fix}(\Delta_c) \subset s'([0, \frac{2}{3}[\setminus \mathbf{Q})$; numbers 0 and $\frac{2}{3}$ are accumulation points of $(s')^{-1}(\text{Fix}(\Delta_c))$.
2. $\text{card}(\text{Fix}(\Delta_c))$ is equal to that of the continuum.



Sturmian words with balanced construction

UPPSALA
UNIVERSITET

Hanna Uscka-Wehlou

Ph.D. Thesis (to be defended in 8 days!):

Digital lines, Sturmian words, and continued fractions

